PLANNC: a projected Lagrangian based implementation for constrained nonlinear network flow problems

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FIRST PART: General overview of the method.

- The Nonlinear Network Flow problem with Side Constraints (NNC).
- The projected Lagrangian methodology.

SECOND PART:

Optimization of the linearized subproblems (NNLC) k : The NOXCB09 package.

- The primal partitioning active-set method.
- Computational results.

THIRD PART:

The use of the NOXCB09 library into a projected Lagrangian methodology

- Structure of the prototype implementation.
- The PLANNC program.
- Preliminary computational results.

CONCLUSIONS

FIRST PART: General overview of the method.

FORMULATION OF THE (NNC) PROBLEM

- (1a) : Objective function: $f: \mathbb{R}^n \to \mathbb{R}$, $f \in \mathcal{C}^2$
- (1b) : Network equations: $A \in \mathbb{R}^{m \times \tilde{n}}$, $r \in \mathbb{R}^m$
- (1c,1d): Nonlinear and linear side constraints: $c: \mathbb{R}^{t_n} \to \mathbb{R}, \ c \in \mathcal{C}^2, \ b_n \in \mathbb{R}^{t_n}$ $T \in \mathbb{R}^{t_l \times \tilde{n}}, \ \operatorname{rang}(T) = t_l, \ b \in \mathbb{R}^{t_l}$

(1e,1f,1g): Arcs and slacks capacities:
$$x, u_x \in \mathbb{R}^n$$

$$z_n, u_{zn} \in \mathbb{R}_n^{\tilde{t}}, \mathbf{I}_n \in \mathbb{R}^{t_n \times \tilde{t}_n}$$

$$z_l, u_{zl} \in \mathbb{R}_l^{\tilde{t}}, \mathbf{I}_n \in \mathbb{R}^{t_l \times \tilde{t}_l}$$

THE PROJECTED LAGRANGIAN METHOD

Algorithm A1: Projected Lagrangian Algorithm for problem (NNC)

• Select some starting values for x^0 , λ^0 and ρ^0 .

MAJOR ITERATION

<u>Do While</u> x^k not optimal for (NNC)

- Linearize the nonlinear side constraints at x^k :
 - $c(x) \approx c^k(x) = c(x^k) + \nabla c(x^k)(x x^k) \tag{2}$
- Minor iteration: Solve the following Nonlinear Network flow problem with Linear side Constraints (NNLC)^k:

$$\operatorname{NNLC}^{k} \left\{ \begin{array}{cccc} \min & \Phi^{k}(x) & (3a) \\ \operatorname{subj. a:} & & \\ & Ax & = r & (3b) \\ & \nabla c(x^{k})x + \mathbf{I}_{n}z_{n} = b_{n}^{k} & (3c) \\ & Tx & + \mathbf{I}_{l}z_{l} = b_{l} & (3d) \\ & 0 \leq x \leq u_{x} & (3e) \\ & 0 \leq z_{n} \leq u_{z_{l}} & (3f) \\ & 0 \leq z_{l} \leq u_{z_{n}} & (3g) \end{array} \right.$$

where:

$$b_n^k = b_n - c(x^k) + \nabla c(x^k) x^k$$

$$\Phi^k(x) = \Phi^k(x; \lambda^k, \dots) \text{ (Merit function)}$$

- Update x^k , λ^k and $\Phi^k(x)$ from $[x^*]^k$ and $[\lambda^*]^k$.
- Set k := k + 1

End Do

RELEVANT FACTS.

Local convergence:

* Robinson (1972) proved the q-quadratically local convergence of the method, if the merit function and the iterates $\{x^{k+1}, \lambda^{k+1}\}$ are taken as:

$$\Phi^{k}(x) = f(x) - \lambda^{k'} (c(x) - c^{k}(x))$$
$$\{x^{k+1}, \lambda^{k+1}\} := \{[x^{*}]^{k}, [\lambda^{*}]^{k}\}$$

Computational implementations:

- MINOS (Murtagh & Saunders 1982).
- Their merit function is a modified augmented Lagrangian function $\Phi^k(x)$:

$$\Phi^{k}(x; x^{k}, \lambda^{k}, \rho^{k}) = f(x) - {\lambda^{k}}' \Big(c(x) - c^{k}(x) \Big) + \frac{\rho^{k}}{2} \|c(x) - c^{k}(x)\|_{2}$$

Efficient solution of the (NNLC)^k subproblems:

- NOXCB09 (Heredia and Nabona 1995).
- Active set method specialized through primal partitioning techniques.

SECOND PART: Optimization of the linearized subproblems (NNLC) k : The NOXCB09 package.

FORMULATION OF THE PROBLEM $(NNLC)^k$

$$(\mathbf{NNLC}) \begin{cases} \min & f(x) & (4a) \\ \text{subj. a} : & Ax = r & (4b) \\ & Tx + \mathbf{I}_z z = b & (4c) \\ & 0 \le x \le u_x & (4d) \\ & 0 \le z \le u_z & (4e) \end{cases}$$

- (4a) : Objective function. $f: \mathbb{R}^n \to \mathbb{R}, \ f \in \mathcal{C}^2$
- (4b) : Network equations $A \in \mathbb{R}^{m \times \tilde{n}}, r \in \mathbb{R}^m$
- (4c) : Linear side constraints $T \in \mathbb{R}^{t \times \tilde{n}}$, rang(T) = t, $b \in \mathbb{R}^t$
- (4d) : Arcs capacities $x, u_x \in \mathbb{R}^n$
- (4e) : Bounds of the slacks $z, u_z \in \mathbb{R}^{\tilde{t}}, \mathbf{I}_z \in \mathbb{R}^{t \times \tilde{t}}$ $(\tilde{t}: number of inequality side constraints)$

Solution procedure for the (NNLC) subproblem.

Algorithm A2: Active set method

Find an initial feasible solution x^0 .

Define the basic \mathcal{B}^0 , superbasic \mathcal{S}^0 and nonbasic \mathcal{N}_0^0 , \mathcal{N}_u^0 variables associated to the current active set at x^0 .

Compute Z^0 , $g_z^0 = Z^{0'}\nabla f(x^0)$ and the Lagrange multipliers of the nonbasic variables σ_0^0 and σ_u^0 .

Set k := 0

<u>Do</u>

Relax some nonbasic variable from its active bound.

Update \mathcal{B}^k , \mathcal{S}^k , \mathcal{N}^k , Z^k and g_z^k .

Do While $||q_z|| > \epsilon_G^k$

Compute a feasible descent direction:

$$H_z^{k}p_z^{k} = -g_z^{k} \quad ; \quad p^k := Z^k p_z^{k}$$

Find the maximal step length: $\bar{\alpha} = \min\{\bar{\alpha}_B, \bar{\alpha}_S\}$.

Linesearch: $\alpha^{*k} = \operatorname{argmin}_{0 < \alpha < \bar{\alpha}} \{ f(x^k + \alpha p^k) \}.$

Update the current solution: $\bar{x}^{k+1} := x^k + \alpha^{*k} p^k$.

$$\underline{\text{If}} \ \alpha^{*k} = \bar{\alpha} \ \underline{\text{then}}$$

Update \mathcal{B}^k , \mathcal{S}^k , \mathcal{N}^k

End If

Update Z^k and q_z^k .

k := k + 1

End Do

Compute the Lagrange multipliers estimates σ_0^k , σ_u^k .

<u>Until</u> $\sigma_0^k \ge -\epsilon_O$ and $\sigma_u^k \le \epsilon_O$

STRUCTURE OF THE PROBLEM (NNLC)

• Equality constraint matrix.

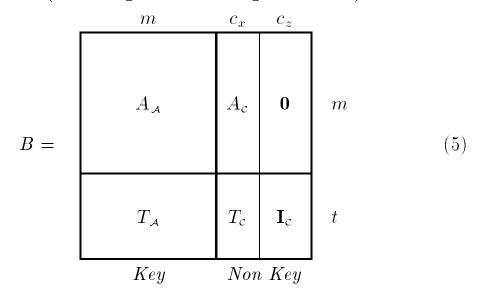
$$M = {}^{m} \left\{ \begin{bmatrix} A & \mathbf{0} \\ T & \mathbf{I}_z \end{bmatrix} = \begin{bmatrix} m & s & n-m-s \\ B & S \end{bmatrix} \quad N \right\}$$

ullet Superbasic S, non-basic N and basic B matrices.

	s_x	s_z				$ \mathcal{N}_x $		$ \mathcal{N}_z $	i
S =	$A_{\mathcal{S}}$	0	m	; N =		$A_{\mathcal{N}}$		0	m
	$T_{\mathcal{S}}$	$\mathbf{I}_{\mathcal{S}}$	t			$T_{\mathcal{N}}$		$\mathbf{I}_{\mathcal{N}}$	t
				m	c_x	c_z			
	1	B =		$A_{\mathcal{A}}$	A_{c}	0	m		
				$T_{\mathcal{A}}$	$T_{\mathcal{C}}$	$\mathbf{I}_{\mathcal{C}}$	t		

PRIMAL PARTITIONING

• Every basis of the (NNLC) problem may be placed in the following form (Kennington and Helgason 1980):



where A_A is a basis (spanning tree) of the network matrix A

- The systems of equations Bv = w and v'B = w can be computed solving:
 - * Systems of equations with the matrix $A_{\mathcal{A}}$, which can be computed efficiently through network flow techniques.
 - * Systems of equations with the general (dense) matrix Q. $Q = [T_c \mid \mathbf{I}_c] T_{\mathcal{A}} A_{\mathcal{A}}^{-1} [A_c \mid \mathbf{0}] \in \mathbb{R}^{t \times t}$

(working basis)

COMPUTATIONS WITH THE BASIC MATRIX B

• Resolution of the linear system Bw = v:

$$w = \begin{bmatrix} w_{\mathcal{A}} \\ - \\ w_{\mathcal{C}} \end{bmatrix} = B^{-1} \begin{bmatrix} v_{\mathcal{A}} \\ - \\ v_{\mathcal{T}} \end{bmatrix} = \begin{bmatrix} A_{\mathcal{A}}^{-1} \left(v_{\mathcal{A}} - [A_{\mathcal{C}} | \mathbf{0}] Q^{-1} \left(v_{\mathcal{T}} - T_{\mathcal{A}} A_{\mathcal{A}}^{-1} v_{\mathcal{A}} \right) \right) \\ Q^{-1} \left(v_{\mathcal{T}} - T_{\mathcal{A}} A_{\mathcal{A}}^{-1} v_{\mathcal{A}} \right)$$
(6a)

$$= \left[\frac{A_{A}^{-1}v_{A} - [\Theta_{c}|\mathbf{0}]Q^{-1}(v_{T} - T_{A}A_{A}^{-1}v_{A})}{Q^{-1}(v_{T} - T_{A}A_{A}^{-1}v_{A})} \right]$$
(6b)

Where $A_{\mathcal{A}}^{-1}A_{\mathcal{C}}=\Theta_{\mathcal{C}}$, is the matrix of the basic cycles of the non-key arcs.

• Algorithm to compute Bw = v:

Algorithm A3:

$$w = B^{-1}v$$
, without $\mathring{\mathcal{C}}$.

$$\boxed{\mathbf{1}} \gamma_1 \stackrel{F.X.}{\longleftarrow} A_{\mathcal{A}}^{-1} v_{\mathcal{A}}$$

$$\mathbf{2} \mid \gamma_2 \longleftarrow v_T - T_A \gamma_1$$

Algorithm A4:

$$w = B^{-1}v$$
, with $\mathring{\mathcal{C}}$.

$$1 \gamma_1 \stackrel{F.X.}{\longleftarrow} A_{\mathcal{A}}^{-1} v_A$$

$$\begin{array}{c|cccc}
\hline
\mathbf{1} & \gamma_1 & \stackrel{F.X.}{\longleftarrow} A_{\mathcal{A}}^{-1} v_A \\
\hline
\mathbf{2} & \gamma_2 & \longleftarrow v_T - T_{\mathcal{A}} \gamma_1 \\
\hline
\mathbf{3} & w_c & \longleftarrow Q^{-1} \gamma_2
\end{array}$$

$$\mathbf{3} \mid w_{c} \longleftarrow Q^{-1} \gamma_{2}$$

$$\boxed{\mathbf{4}} \quad w_{\mathcal{A}} \stackrel{\text{\tiny A4.1}}{\longleftarrow} \gamma_1 - \Theta_{\mathcal{C}} w_{\mathcal{C}_x}$$

Product
$$w = Zv = \begin{bmatrix} -B^{-1}S \\ \hline I \\ \hline 0 \end{bmatrix}v$$

• Option A: first $\gamma_A = -Sv$ and then $w = B^{-1}\gamma$

$$Sv = \begin{bmatrix} A_{\mathcal{S}} & \mathbf{0} \\ T_{\mathcal{S}} & \mathbf{I}_{\mathcal{S}} \end{bmatrix} \begin{bmatrix} v_x \\ v_z \end{bmatrix} = \begin{bmatrix} A_{\mathcal{S}}v_x \\ \hline T_{\mathcal{S}}v_x + \mathbf{I}_{\mathcal{S}}v_z \end{bmatrix}$$
(7)

Algorithm A5 : $-B^{-1}(Sv)$

$$\boxed{1} \quad \gamma_A \longleftarrow -A_{\mathcal{S}} v_x$$

$$\boxed{2} \gamma_T \longleftarrow -T_S v_x - \mathbf{I}_S v_z$$

$$3 \quad w_{\mathcal{B}} \stackrel{\text{\tiny A3}}{\longleftarrow} B^{-1} \gamma$$

• Option B: to develop $B^{-1}S$.

$$w_{\mathcal{B}} = \begin{bmatrix} w_{\mathcal{A}} \\ w_{\mathcal{C}} \end{bmatrix} = -B^{-1}Sv = \begin{bmatrix} -\Theta_{\mathcal{S}}v_x - [\Theta_{\mathcal{C}}|\mathbf{0}]Q^{-1}\left(T_{\mathcal{A}}\Theta_{\mathcal{S}}v_x - T_{\mathcal{S}}v_x - \mathbf{I}_{\mathcal{S}}v_z\right) \\ Q^{-1}\left(T_{\mathcal{A}}\Theta_{\mathcal{S}}v_x - T_{\mathcal{S}}v_x - \mathbf{I}_{\mathcal{S}}v_z\right) \end{bmatrix}$$
(8)

Algorithm A6 : $w_{B} = Zv$.

$$1 \mid \gamma_1 \stackrel{\text{\tiny A4.2}}{\longleftarrow} \Theta_{\scriptscriptstyle \mathcal{S}} v_x$$

$$\boxed{2} \quad \gamma_2 \longleftarrow T_{\mathcal{A}} \gamma_1 - T_{\mathcal{S}} v_x - \mathbf{I}_{\mathcal{S}} v_z$$

$$\begin{array}{|c|c|c|}\hline \mathbf{1} & \gamma_1 \xleftarrow{^{\mathrm{A4.2}}} \Theta_{\mathcal{S}} v_x \\ \hline \mathbf{2} & \gamma_2 \longleftarrow T_{\mathcal{A}} \gamma_1 - T_{\mathcal{S}} v_x - \mathbf{I}_{\mathcal{S}} v_z \\ \hline \mathbf{3} & [w'_{\mathcal{C}_x} \mid w'_{\mathcal{C}_z}] \longleftarrow -Q^{-1} \gamma_2 \\ \hline \mathbf{4} & w_{\mathcal{A}} \xleftarrow{^{\mathrm{A4.1}}} \gamma_1 - \Theta_{\mathcal{C}} w_{\mathcal{C}_x} \\ \end{array}$$

$$4 \quad w_{\mathcal{A}} \stackrel{\text{\tiny A4.1}}{\longleftarrow} \gamma_1 - \Theta_{\mathcal{C}} w_{\mathcal{C}_x}$$

OTHER IMPLEMENTATION ISSUES

- Linesearch:
 - * Bertsekas method (Bertsekas (1982), Toint & Tuyttens (1991)).
 - * Bactracking linesearch with quadratic and cubic fit.
- Descent direction computation: $\tilde{H}_z p_z = -g_z(x)$
 - * (TNM) Truncated Newton Method (Dembo & Steihaug (1983), Toint & Tuyttens (1991)).
 - * (QNM) Quasi-Newton method (Murtagh & Saunders (1978)).
- Working basis update $[Q^{-1}]^k$: cada cert nombre d'iteracions
 - * Product form of the inverse :

$$[Q^{-1}]^k = \prod_{t \ge i \ge 1} E_r^i$$

- $ightharpoonup P^3$ algorithm of Hellerman & Rarick : $[Q^{-1}]^k = R\left(\prod_{t \geq i \geq 1} \tilde{E}_r^i\right) P$
- ▶ Partial pivoting.
- * LU factorization.
- ▶ Subroutines F01BRF and F04AXF (NAG library).

EIO/UPC problem tests

- 110 problemes (NNLC).
 - * DIMACS problems:
 - ${\,\vartriangleright\,}$ Random generators of linear network flow problems (FL) :

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Rmfgen (Goldfarb, Grigoriadis). Grid-on-Torus (Goldberg).
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Gridgen (Lee & Orlin).

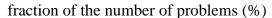
- \triangleright Addition of side constraints to the DIMACS problems: Di2no.
- ▶ Objective functions :

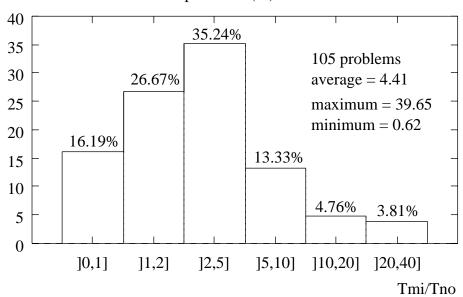
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Namur (Toint & Tuyttens). EIO1.
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- * Real Short term Hydro-Thermal Coordination problems.
- Characteristics of the EIO/UPC test collection:
 - * 1524 arcs, 360 nodes \rightarrow 23832 arcs and 6589 nodes.
 - * Ratio (# of side const./ # network equations) : $1\% \rightarrow 100\%$
 - * # of nonzero elements of the S.C. : $0.02\% \rightarrow 10\%$
 - * # of superbasic variables at the optimal solution: $0 \rightarrow 3352$

COMPUTATIONAL RESULTS (summary)

• Efficiency of NOXCB09 vs. MINOS 5.3:





• Dependency of the efficiency on the size of the problem:

	Max	imal d	dimensions	$ m t_{MI}/t_{NO}$				
Scale	arcs	nodes	# S.C.	Average	Maximum	Minimum		
Large	18000	3000	750	5.10	17.96	0.84		
Medium	8064	2479	840	2.87	10.00	0.92		
Small	2256	697	240	1.34	2.47	0.62		

* The efficiency of NOXCB09 seems to increase with the dimension of the (NNLC) problem.

THIRD PART:

The use of the NOXCB09 library into a projected Lagrangian methodology

STRUCTURE OF THE IMPLEMENTATION

Algorithm A7: a NOXCB09 based Projected Lagrangian method

Select some starting values for x^0 , λ^0 and ρ^0 .

Set k := 0

 $\underline{\text{Do}}$

Linearize the nonlinear side constraints at x^k .

<u>Do</u>

Solve $(\mathbf{NNLC})^k$ with NOXCB09 and the merit function $\Phi^k(x; x^k, \lambda^k, \rho^k)$. This optimization provides $[x^*]^k$ and $[\lambda^*]^k$.

 $\underline{\text{If}} (\mathbf{NNLC})^k \text{infeasible } \underline{\text{then}}$

Relax the RHS of the linearized n.s.c. b_n^k .

End If

 $(NNLC)^k$ feasible Until

Find new values of x^k and λ^k from $[x^*]^k$ and $[\lambda^*]^k$.

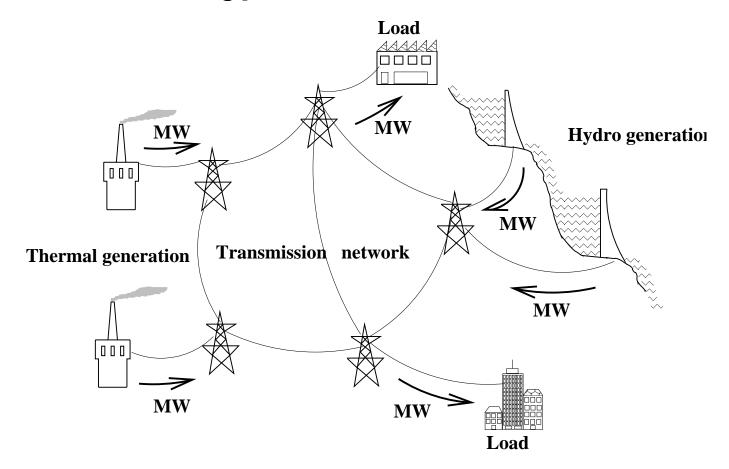
Update de penalty parameter ρ^k , if necessary.

Set k := k + 1

 x^k optimal for (NNC) Until

• Updates rules the update rules follows those suggested by Murtagh & Saunders (1982) and implemented in MINOS.

- PLANNC, a projected Lagrangian NOXCB09 based package to solve the (NNC) problem.
- Applied to solve the **Short-Term Hydrothermal Scheduling** problem:



PRELIMINARY COMPUTATIONAL RESULTS

Table I: characteristics of the **(NNC)** problems.

Problem	-	Power	syster	n size		(NNC) problem size				
${\bf ident.}$	Nr	Nu	Nm	Nb	Ni	arcs	nodes	n.s.c.	l.s.c.	
A24x	3	4	6	5	24	888	290	24	96	
B24x	6	4	6	5	24	1152	355	24	96	
B48x	6	4	6	5	48	2256	697	48	192	

 ${\bf Table~II}: {\bf Computational~results}.$

Problem	Major/N	Inor iter	C.	PU (s) 1		$ m millisec/iter^1$		
ident.	PLANNC	MINOS	PLANNC	MINOS	ratio	PLANNC	MINOS	ratio
A24x	3/2384	53/2030	2.3	4.8	2.08	0.96	2.3	2.4
B24x	3/3947	99/3855	3.1	15.3	4.93	0.78	3.9	5.0
B48x	3/7988	119/4656	6.4	36.2	5.65	0.80	7.7	9.6

¹ SUN UltraSparc 2

CONCLUSIONS

- A projected Lagrangian methodology for the (NNC) problem has been proposed.
- This methodology explotes the network structure of the problem applying the NOXCB09 package to solve the linearized subproblems $(\mathbf{NNLC})^k$. The efficiency of this code has been proved through extensive computational tests.
- A prototype implementation, (PLANNC), of the proposed specialized projected Lagrangian method, has been implemented and used to solve a limited set of (NNC) problems.
- The preliminary computational results seems to corroborate the efficiency of the method, although more computational evidence is needed. These results also reveals the posibility of increasing the convergence rate with a more careful implementation.